Objectives

You should be able to ...

Lambda Calculus

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The purposes of this lecture is to introduce lambda calculus and explain the role it has in programming languages.

- **Explain** the three constructs of λ -calculus.
- Given a syntax tree diagram, write down the equivalent λ -calculus term.
- ▶ Perform a beta-reduction.

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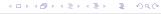
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3. Functions (Also called abstractions.)

 $\lambda x.x$ $\lambda ab.fab$ $\lambda xy.g(\lambda z.zf)yx$



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- ▶ Used extensively in research. The "little white mouse" of computer science.
- ▶ We can implement this trivially in Haskell.

$$\lambda x.x = \langle x \rangle - x.$$

Examples

 $\lambda x.x$ "The identity"

 $\lambda x.xx$ "Delta"

 λ ab.fabxy

 $(\lambda ab.fab)$ xy

 $(\lambda a.\lambda b.fab)$ xy

 $(\lambda f x.x f)(\lambda g.g x)(\lambda f.f)z y$

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Syntax Trees

Example 1	Example 2	Example 3	Example 4
·	λγ	@ @ z	λz \downarrow \Diamond λx λy
$\lambda x \downarrow x$		λy \times \downarrow y	@ @ x z y z
$\lambda x.x$	λy.yx	$(\lambda y.y)xz$	$\lambda z.(\lambda x.xz)(\lambda y.yz)$



Function Application

Objectives

$$(\lambda x.M)N \mapsto [N/x]M$$

$$[N/x] y = y$$

$$[N/x] x = N$$

$$[N/x] (M P) = ([N/x]M [N/x]P)$$

$$[N/x] (\lambda y.M) = \lambda y.[N/x]M$$

$$[N/x] (\lambda x.M) = \lambda x.M$$



Bound and Free

- ▶ The λ creates a binding.
- ► An occurance of the the variable inside the function body is said to be *bound*.
- ▶ Bound variables occur "under the λ " that binds them.

Examples: Where are the free variables? To which lambdas are bound variables bound?

