Dr. Mattox Beckman

University of Illinois at Urbana-Champaign Department of Computer Science

•00000

- ▶ Show how to build complex parsers by composing simpler parsers.
- Use monads to hide the mechanics of plumbing the input.
- Build a small parser library similar to the Parsec combinator parser library in HASKELL.

The Problem

- ► Recursive descent parsers are easy to write.
 - But plumbing the input is a bit tedious.
 - ▶ And sometimes the common prefix problem is a real problem.
 - And we can't really compose them.
- ► So we'll build a parser *combinator* library instead.

A Parser

- We begin by defining a type.
- ► The newtype is like data but with only one constructor.
 - Compiler can handle this more efficiently.
- ▶ The run function unboxes a parser so we can run it.

```
newtype Parser t = Parser (String -> [(t,String)])
zrun (Parser p) = p
```

Our First Parser: Parsing a Character

The char Parser

- Single quotes are for single characters.
- Double quotes are for strings (lists of characters).

```
Main> run (char 'a') "asdf"
[('a',"sdf")]
Main> run (char 'a') "qwert"
[]
```

Predicates and Parsers

- oneOf takes a list of characters and succeeds if the input is one of them.
- ▶ In real life you might want to build a lookup table.

```
1 \text{ oneOf } xx =
   Parser (\inp -> case inp of
                       (s:ss) \mid s \cdot elem \times xx \rightarrow [(s.ss)]
3
                       otherwise
                                                -> [])
5
6 digit = oneOf ['0'..'9']
Main> run (oneOf "asb") "sb"
[('s',"b")]
Main> run (oneOf "asb") "xsb"
Г٦
Main> run digit "42"
[('4',"2")]
```

Making It a Higher Order Function

- **sat** takes a predicate that it can run on the character.
- ► Compare with oneOf.

```
1 \text{ one } 0 \text{ f } xx =
    Parser (\inp -> case inp of
                          (s:ss) \mid s \cdot elem \times xx \rightarrow [(s,ss)]
                          otherwise
                                                       -> [])
6 sat pred =
    Parser (\inp -> case inp of
                          (s:ss) \mid pred s \rightarrow [(s,ss)]
                          otherwise
                                                    -> [])
10
11 \text{ digit} = \text{sat} (\x -> x >= '0' \&\& x <= '9')
```

Adding a Choice Operator

- ▶ We want to compose two parsers together.
- ▶ If the first fails, we will try the second.

```
1 (Parser p1) <|> (Parser p2) =
2  Parser (\inp -> take 1 $ p1 inp ++ p2 inp)

Main> run (digit <|> (char 'a')) "12ab"
[('1',"2ab")]

Main> run (digit <|> (char 'a')) "a2ab"
[('a',"2ab")]

Main> run (digit <|> (char 'a')) "xa2ab"
[]
```

Recursion

Come and see the plumbing inherent in the system!

```
rstring [] = Parser (\inp -> [([],inp)])
2 rstring (s:ss) = Parser (\inp ->
     case run (char s) inp of
       [(c,r1)] -> case run (rstring ss) r1 of
                     [(cs,rr)] -> [(c:cs,rr)]
                    _ -> []
                -> [])
> run (rstring "Arthur Dent") "Arthur Dent"
[("Arthur Dent","")]
```

- We have created a parser using recursion, but this is painful.
- ► What operation do you know that unpacks a data structure, propagates success cases, and aborts computation after a failure?



Enter the Monad - Functor

```
instance Functor Parser where
   fmap f (Parser p1) =
       Parser (\inp -> [(f t, s) |
                          (t,s) \leftarrow p1 inp]
6 sdi :: Parser Integer
7 sdi = Parser (\inp -> case run digit inp of
                           [(d, dd)] -> [(read [d], dd)]
                          otherwise -> [])
 sdi = "single digit integer," not "strategic defense initiative"
Main> run sdi "123"
[(1,"23")]
Main> run (fmap (+1) sdi) "123"
[(2,"23")]
```

Enter the Monad - Applicative

```
pinstance Applicative Parser where
pure a = Parser (\inp -> [(a,inp)])
(Parser p1) <*> (Parser p2) =
Parser (\inp -> [(v1 v2, ss2) |
(v1,ss1) <- p1 inp,
(v2,ss2) <- p2 ss1])

Main> run ( (+) <$> sdi <*> sdi) "456"
[(9,"6")]
```

Enter the Monad

Remember that f takes data from the first parser and returns a new parser.

```
instance Monad Parser where
   (Parser p) >>= f =
       Parser (\inp -> concat [run (f v) inp'
                               | (v,inp') <- p inp])</pre>
Main> run (sdi >>= (\x -> sdi >>= (\y -> return x + y))
          "8675309"
[(14, "75309")]
Main> run (do x < - sdi
              y <- sdi
              return x + y "123"
[(3,"3")]
```

Recursion, Revisited

▶ Using do notation, we can really clean up our code.

Before

```
rstring [] = Parser (\inp -> [([],inp)])

rstring (s:ss) = Parser (\inp ->
case run (char s) inp of

[(c,r1)] -> case run (rstring ss) r1 of
[(cs,rr)] -> [(c:cs,rr)]
-> []
-> []
```

Recursion, Revisited

▶ Using do notation, we can really clean up our code.

After

Many and Many1

```
nmany p = next <|> return ""
    where next = do v <- p
                      vv <- many p
                      return (v:vv)
6 \text{ many 1} p = do v < - p
               vv <- many p
                return (v:vv)
10 spaces = many (oneOf " ")
```

Returning a Type

```
1 data Exp = IntExp Integer
          OpExp String Exp Exp
      deriving Show
5 int :: Parser Exp
6 int = do digits <- many1 digit
          spaces
          return (IntExp $ read digits)
Main> run int "1234 567"
[(IntExp 1234, "567")]
Main> run (many int) "10 20 30 40"
[([IntExp 10,IntExp 20,IntExp 30,IntExp 40],"")]
```

Operators

```
oper o = do v <- string o
             spaces
             return $ OpExp v
4 chainl1 p op = p >>= rest
    where rest x = do o \leftarrow op
                      q -> v
6
                       rest (o x v)
                   <l> return x
9 expr = chainl1 term (oper "+")
10 term = int <|> parens expr
Main> run expr "10 + 20 + 30"
[(OpExp "+" (OpExp "+" (IntExp 10) (IntExp 20)) (IntExp 30),"")]
Main> run expr "10 + (20 + 30)"
[(OpExp "+" (IntExp 10) (OpExp "+" (IntExp 20) (IntExp 30)),"")]
```

Longer Example

Longer Example, II

- ► The try allows for backtracking.
- ▶ There are many packages: parsec, attoparsec, and megaparsec.