

# Objects

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## Objectives

You should be able to ...

In this lecture, we extend the idea of local state from last time to create a simple implementation of objects and discuss its limitations. We will also show the message dispatch model of objects, which allows for inheritance and virtual functions.

Your objectives:

- ▶ Be able to explain what an object is.
- ▶ Implement an object using records and HOFs.
- ▶ Implement an object using a message dispatcher.

## Preliminaries

- ▶ We will use the following functions during our discussion:

```
1 let pi1 (x,y) = x
2 let pi2 (x,y) = y
3 let report (x,y) = print_string "Point: ";
4     print_int x;
5     print_string ",";
6     print_int y;
7     print_newline ()
8 let movept (x,y) (dx,dy) = (x+dx,y+dy)
```

## Point

Here is an example of a point using local state.

```
1 let mktPoint myloc =
2   let myloc = ref myloc in
3     ( myloc,
4       (fun () -> pi1 !myloc),
5       (fun () -> pi2 !myloc),
6       (fun () -> report !myloc),
7       (fun dl -> myloc := movept !myloc dl) )
```

- ▶ This defines a tuple of functions that share a common state.
- ▶ It is cumbersome to use.

```
let (lref,getx,gety,show,move) = mktPoint (2,4);;
```

## Improvement: Use Records

```

1 type point = {
2   loc : (int * int) ref; getx : unit -> int;
3   gety : unit -> int; draw : unit -> unit;
4   move : int * int -> unit;
5 }
6 let mkrPoint newloc =
7   let myloc = ref newloc in
8   { loc = myloc;
9     getx = (fun () -> pi1 !myloc);
10    gety = (fun () -> pi2 !myloc);
11    draw = (fun () -> report !myloc);
12    move = (fun dl -> myloc := movept !myloc dl)}

```



## Message Dispatching

Last time we said that an object is a kind of data that can *receive messages* from the program or other objects.

- ▶ Q: How do we normally represent messages?
- ▶ A: With strings!

Let a point object be a function that takes a string and returns an appropriate function matching that string.

- ▶ Question: Suppose `p` is our point object. What will be its type?



## Adding Self

By the way, this lecture is really about recursion.

```

1 let mkPoint newloc =
2   let rec this =
3     { loc = ref newloc;
4       getx = (fun () -> pi1 !(this.loc));
5       gety = (fun () -> pi2 !(this.loc));
6       draw = (fun () -> report !(this.loc));
7       move = (fun dl ->
8                 this.loc := movept !(this.loc) dl)}
9   in this;;

```

We can store "this" explicitly in the record if we want.



## mkPoint

```

1 let mkPoint x y =
2   let x = ref x in
3   let y = ref y in
4   fun st ->
5     match st with
6     | "getx" -> (fun _ -> !x)
7     | "gety" -> (fun _ -> !y)
8     | "movx" -> (fun nx -> x := !x + nx; nx)
9     | "movy" -> (fun ny -> y := !y + ny; ny)
10    | _ -> raise (Failure "Unknown message.")

```

All methods now have to have type `int -> int`.



## Subclassing

- ▶ Warmup exercise: How would we add a `report` method?
- ▶ Another one: How would we add `this` support?

Let's say we want a `fastpoint`, which moves twice as fast as the original point. What does it mean for `fastpoint` to be a *subclass* of `point`?

- ▶ `fastpoint` should respond to the same messages.
  - ▶ It may override some of them.
  - ▶ It may add its own.
  - ▶ It may **not** remove any methods.
- ▶ The `fastpoint` object will need access to some of the data in `point`.



## Code: point

```

1 let mkSuperPoint x y =
2   let x = ref x in
3   let y = ref y in
4   ((x,y), (* This part returns the local state *))
5   fun st ->
6     match st with
7     | "getx" -> (fun _ -> !x)
8     | "gety" -> (fun _ -> !y)
9     | "movx" -> (fun nx -> x := !x + nx; nx)
10    | "movy" -> (fun ny -> y := !y + ny; ny)
11    | _ -> raise (Failure "Unknown message.");;
12 val mkSuperPoint : int -> int ->
13   (int ref * int ref) * (string -> int -> int) = <fun>

```



## Implementing

- ▶ Two entities involved: the superclass (`point`) and the subclass (`fastpoint`)
- ▶ `fastpoint` needs to create an instance of `point`.
- ▶ `point` construction needs to return the "public" data to `fastpoint`.
- ▶ `fastpoint` returns a dispatcher:
  - ▶ If the `fastpoint` dispatcher can handle a message, it does.
  - ▶ Otherwise, it *sends the message* to `point`.



## Code: fastpoint

```

1 let mkFastpoint x y =
2   let ((x,y),super) = mkSuperPoint x y in
3   fun st ->
4     match st with
5     | "movx" -> (fun nx -> x := !x + 2 * nx; nx)
6     | "movy" -> (fun ny -> y := !y + 2 * ny; ny)
7     | _ -> super st;;

```

- ▶ This technique is flexible; we can add methods very easily.
- ▶ But it's also slow. Imagine if we had a chain of 20 classes!



# C++

- ▶ Methods and variables are kept in a table: a fixed location.
- ▶ "this" is an implicit argument, allowing only one copy of the function to be needed.
- ▶ Virtual methods are kept in a *vtable*, which counts as local data.

Local data for point or fastpoint:

x	value of x
y	value of y
vtable	pointer to vtable

Vtable for point:

movx	pointer to point.movx
movy	pointer to point.movy

(fastpoint vtable is similar.) getx, etc. is static.

# Discussion

- ▶ Other languages (i.e., smalltalk) use a technique very similar to this one.
- ▶ Java uses the "every object is of type Object" technique.
- ▶ A strong type system makes it somewhat cumbersome to simulate objects. You either have to:
  - ▶ define a new type to encompass all objects, or
  - ▶ force all methods to have the same type.
- ▶ Important concept: *polymorphism* — when functions can operate on multiple types. (This is different than *overloading* — when multiple functions exist with the same name, but different inputs.)