

Loop Equations

► We need to solve five equations.

```

{q}
Si
{inv : p} {bd : t}
while B do
  {p ∧ B}
  S
  {p}
od
{p ∧ ¬B}
{r}
    
```

1. $\{q\}S_i\{p\}$
2. $\{p \wedge B\}S\{p\}$
3. $p \wedge \neg B \rightarrow r$
4. $p \rightarrow t \geq 0$
5. $\{p \wedge B \wedge t = z\}S\{t < z\}$



Example 1 – Partial Correctness

Example 1

```

s := 0;
i := 0;
while (i < |A|) do
  s := s + A[i];
  i := i + 1
od
    
```

What are these equations?

- $\{q\}S_i\{p\}$
- $\{p \wedge B\}S\{p\}$
- $p \wedge \neg B \rightarrow r$

Solutions:

- $\{\mathbf{true}\}s := 0; i := 0\{i \leq |A| \wedge s = \sum_0^{i-1} A[i]\}$
- $\{i \leq |A| \wedge s = \sum_0^{i-1} A[i] \wedge i < |A|\}S\{i \leq |A| \wedge s = \sum_0^{i-1} A[i]\}$
- $i \leq |A| \wedge s = \sum_0^{i-1} A[i] \wedge i \geq |A| \rightarrow s = \sum_0^{|A|-1} A[i]$



Example 2 – Partial Correctness

Example 2

```

while (a > 0) do
  a, b := b mod a, a
od
    
```

What are these equations?

- $\{q\}S_i\{p\}$
- $\{p \wedge B\}S\{p\}$
- $p \wedge \neg B \rightarrow r$

Solutions:

- No initialization!
- $\{gcd(a, b) = gcd(a', b') \wedge a > 0\}S\{gcd(a, b) = gcd(a', b')\}$
- $gcd(a, b) = gcd(a', b') \wedge a = 0 \rightarrow b = gcd(a', b')$



How to Pick a Loop Invariant

- The loop invariant is a weaker version of the postcondition.
- $p \wedge \neg B \rightarrow r$
- The loop's job is to incrementally make B false.
- So, to pick a loop invariant, you need to weaken the postcondition.

Ways to Weaken

- Replace a constant with a range.
- Add a disjunct.
- Remove a conjunct.



Example 3

$$|f(x)| < \varepsilon \wedge \delta < \varepsilon$$

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$$|f(x)| < \varepsilon$$

Making Progress

- ▶ What does it mean to “make progress toward termination?”
- ▶ Consider a function on integers ...
- ▶ A function on lists ...
- ▶ A function on Hydras ...

The Total Correctness Formulas

- ▶ $p \rightarrow t \geq 0$
- ▶ $\{p \wedge B \wedge t = z\}S\{t < z\}$

Example 1 – Total Correctness

Example 1

```
s := 0;
i := 0;
while (i < |A|) do
  s := s + A[i];
  i := i + 1
od
```

What are these equations?

- ▶ $p \rightarrow t \geq 0$
- ▶ $\{p \wedge B \wedge t = z\}S\{t < z\}$

Solution:

- ▶ $i \leq |A| \wedge s = \sum_0^{i-1} A[i] \rightarrow t \geq 0$
- ▶ $\{i \leq |A| \wedge s = \sum_0^{i-1} A[i] \wedge i < |A| \wedge t = z\}S\{t < z\}$
- ▶ Let $t = |A| - i$.



Example 2 – Total Correctness

Example 2

```
while (a > 0) do
  a, b := b mod a, a
od
```

What are these equations?

- ▶ $p \rightarrow t \geq 0$
- ▶ $\{p \wedge B \wedge t = z\}S\{t < z\}$

Solutions:

- ▶ $a > 0 \rightarrow t \geq 0$
- ▶ (Too big to fit. But notice a always decreases!)

