

Objectives

You should be able to...

Prolog

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In this lecture, we will introduce Prolog.

- ▶ Explain how Prolog uses a unification to drive computation.
- ▶ Write some simple programs in Prolog.

Logic

Question: How do you decide truth?

- ▶ Start with some *objects*.
"socrates," "john," "mary"
- ▶ Write down some *facts* (true statements) about those objects.
 - ▶ Facts express either properties of the object, or
"socrates is human"
 - ▶ relationship to other objects.
"mary likes john"
- ▶ Write down some *rules* (facts that are true if other facts are true).
"if X is human then X is mortal"
- ▶ Facts and rules can become *predicates*.
"is socrates mortal?"

First-Order Predicate Logic

First-order predicate logic is one system for encoding these kinds of questions.

- ▶ Predicate means that we have functions that take objects and return "true" or "false."
`human(socrates)` .
- ▶ Logic means that we have *connectives* like and, or, not, and implication.
- ▶ First order means that we have variables (created by "for all" and "there exists"), but that they only work on objects.
 $\forall X . \text{human}(X) \rightarrow \text{mortal}(X)$.

History

- ▶ Starting point: First-order predicate logic.
- ▶ Realization: computers can reason with this kind of logic.
- ▶ Impetus was the study of *mechanical theorem proving*
- ▶ Developed in 1970 by Alain Colmerauer and Rober Kowalski and others
- ▶ Uses: databases, expert systems, AI

The Two Questions

What is the nature of data?

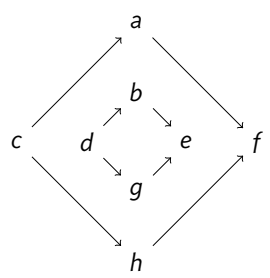
Prolog data consists of *facts* about *objects* and logical *rules*.

What is the nature of a program?

A program in Prolog is a set of facts and rules, followed by a *query*.



The Database



- 1 **human**(socrates).
- 2 **fatherof**(socrates, jane).
- 3 **fatherof**(zeus, apollo).

- 1 **connected**(c, a).
- 2 **connected**(c, h).
- 3 **connected**(d, b).
- 4 **connected**(d, g).
- 5 **connected**(a, f).
- 6 **connected**(h, f).
- 7 **connected**(b, e).
- 8 **connected**(g, e).

Rules

- 1 **mortal**(X) :- **human**(X).
- 2 **human**(Y) :- **fatherof**(X, Y), **human**(X).
- 3
- 4 **pathfrom**(X, Y) :- **connected**(X, Y).
- 5 **pathfrom**(X, Y) :- **connected**(X, Z),
- 6 **pathfrom**(Z, Y).

- ▶ Capital letters are variables.
 - ▶ Appearing left of :- means "for all"
 - ▶ Appearing right of :- means "there exists"

$$\forall x. human(x) \rightarrow mortal(x).$$

$$\forall y. (\exists x. fatherof(x, y) \wedge human(x)) \rightarrow human(y)$$



How It Works

Programs are executed by searching the database and attempting to perform unification.

```
1?- human(socrates).    -- listed, therefore true
2?- mortal(socrates).  -- not listed
```

Relevant rules:

```
1 human(socrates).
2 human(Y) :- fatherof(X,Y), human(X).
3 mortal(X) :- human(X).
```

Socrates is not listed as being mortal, but mortal(socrates) unifies with mortal(X) if we replace X with socrates. This gives us a *subgoal*. Replace X with socrates and try it...



Another Example

```
▶ ?- mortal(jane).
```



How It Works, Next Step

Replace X with socrates in this rule:

```
1 mortal(X) :- human(X).
```

to get

```
1 mortal(socrates) :- human(socrates).
```

Since human(socrates) is in the database, we know that mortal(socrates) is also true.



Another Example

```
▶ ?- mortal(jane).
    mortal(X) :- human(X).
```



Another Example

```

▶ ?- mortal(jane).
   mortal(jane) :- human(jane).

```



Another Example

```

▶ ?- mortal(jane).
   mortal(jane) :- human(jane).
     ▶ human(jane)
       human(Y) :- fatherof(X,Y), human(X).

```



Another Example

```

▶ ?- mortal(jane).
   mortal(jane) :- human(jane).
     ▶ human(jane)

```



Another Example

```

▶ ?- mortal(jane).
   mortal(jane) :- human(jane).
     ▶ human(jane)
       human(jane) :- fatherof(X,jane), human(X).

```



Another Example

- ▶ `?- mortal(jane).`
- `mortal(jane) :- human(jane).`
- ▶ `human(jane)`
- `human(jane) :- fatherof(X,jane), human(X).`
- ▶ `fatherof(X,jane)`



Another Example

- ▶ `?- mortal(jane).`
- `mortal(jane) :- human(jane).`
- ▶ `human(jane)`
- `human(jane) :- fatherof(socrates,jane), human(socrates).`
- ▶ `fatherof(X,jane)`
- ▶ `fatherof(socrates,jane)`



Another Example

- ▶ `?- mortal(jane).`
- `mortal(jane) :- human(jane).`
- ▶ `human(jane)`
- `human(jane) :- fatherof(X,jane), human(X).`
- ▶ `fatherof(X,jane)`
- ▶ `fatherof(socrates,jane)`



Another Example

- ▶ `?- mortal(jane).`
- `mortal(jane) :- human(jane).`
- ▶ `human(jane)`
- `human(jane) :- fatherof(socrates,jane), human(socrates).`
- ▶ `fatherof(X,jane)`
- ▶ `fatherof(socrates,jane)`
- ▶ `human(socrates)`



You Try ...

- ▶ Given the connected rules, try to come up with a predicate `exactlybetween(A,B,C)` that is true when B is connected to both A and C.
- ▶ Now make a predicate `between(A,B,C)` that is true if there's a path from A to B to C.

```

1 exactlybetween(A,B,C) :- connected(A,B), connected(B,C).
2
3 between(A,B,C) :- pathfrom(A,B), pathfrom(B,C).

```



More Than Just Yes or No

- ▶ Prolog can also give you a list of elements that make a predicate true. Remember unification.

```

1 ?- fatherof(Who,apollo).
2 Who = zeus
3
4 ?- pathfrom(c,X).
5 X = a ;
6 X = h ;
7 X = f ;
8 X = f ;
9 No

```

The semicolon is entered by the user — it means to keep searching.



Tracing pathfrom

```

1 ?- pathfrom(c,X).
2 ----> pathfrom(c,Y) :- connected(c,Y).
3 X = a ;

```

When we hit semicolon, we tell it to keep searching. So we *backtrack* through our database to try again.

```

1 pathfrom(c,Y) :- connected(c,Y).
2 ----> X = h ;

```

We tell it to try again with this one, too. At this point, we no longer have any rules that say that c is connected to something.



Tracing pathfrom, II

```
1 pathfrom(c,Y) :- connected(c,Z), pathfrom(Z,Y).
```

We will first find something in the database that says that *c* is connected to some *Z*, and then check if there is a path between *Z* and *Y*.

We find *a* and *h* as last time. When we check *a*, we check for `pathfrom(a,Y)`, and find that `connected(a,f)` is in the database. The same thing happens for *h*, which is why *f* is reported as an answer twice.



Arithmetic via the is Keyword.

```
1 fact(0,1).
2 fact(N,X) :- M is N-1, fact(M,Y), X is Y * N.
3 ?- fact(5,X).
```

- ▶ Unify `fact(5,X)` with `fact(N,X)`.
`fact(5,X) :- M is 5-1, fact(M,Y), X is Y * 5.`
- ▶ Next compute *M*.
`fact(5,X) :- 4 is 5-1, fact(4,Y), X is Y * 5.`
- ▶ Recursive call sets *Y* to 24.
`fact(5,X) :- 4 is 5-1, fact(4,24), X is 24 * 5.`
- ▶ Compute *X*.
`fact(5,120) :- 4 is 5-1, fact(4,24), 120 is 24 * 5.`



Lists

- ▶ Empty list: `[]`
- ▶ Singleton list: `[x]`
- ▶ List with multiple elements: `[x,y,[a,b],c]`
- ▶ Head and tail representation: `[H|T]`

Differences:

- ▶ Prolog lists are *not* monotonic!



List Example: mylength

The `length` predicate is built in.

```
1 mylength([],0).
2 mylength([_|_],X) :- mylength(T,Y),
3                       X is Y + 1.
4
5 ?- mylength([2,3,4,5],X).
6 X = 4 ;
7 No
```



List Example: Sum List

```

1 sumlist([],0).
2 sumlist([H|T],X) :- sumlist(T,Y),
3                       X is Y + H.
4
5 ?- sumlist([2,3,4,5],X).
6 X = 14

```

Try writing list product now!



List Example: Append

```

1 myappend([],X,X).
2 myappend([H|T],X,[H|Z]) :- myappend(T,X,Z).
3 ?- myappend([2,3,4],[5,6,7],X).
4 X = [2, 3, 4, 5, 6, 7] ;
5 No
6 ?- myappend(X,[2,3],[1,2,3,4]).
7 No
8 ?- myappend(X,[2,3],[1,2,3]).
9 X = [1] ;
10 No

```



List Example: Reverse

Accumulator recursion works in Prolog, too!

```

1 myreverse(X,Y) :- aux(X,Y, []).
2 aux([],Y,Y).
3 aux([HX|TX],Y,Z) :- aux(TX,Y,[HX|Z]).
4 ?- myreverse([2,3,4],Y).
5 Y = [4, 3, 2]

```

$\text{myreverse}([2,3,4],Y) \rightarrow \text{aux}([2,3,4],Y,[]) \rightarrow \text{aux}([3,4],Y,[2]) \rightarrow$
 $\text{aux}([4],Y,[3,2]) \rightarrow \text{aux}([],Y,[4,3,2]) \rightarrow \text{aux}([], [4,3,2], [4,3,2]) \rightarrow$
 $\text{myreverse}([2,3,4], [4,3,2])$

